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Ada COMPILER
VALIDATION SUMMARY REPORT:
Certificate Number: 921210W1.11302
Alsys, Inc.
AlsyCOMP 019, Version 5.3.1
CompuAdd 433 under MS/DOS v5.0 and Phar Lap v4.0 =>
Intel iSBC 186/100

(FINAL)

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Certificate Information

The following Ada implementation was tested and determined to pass ACVC 1.11. Testing was completed on 10 December 1992.

Compiler Name and Version: AlsyCOMP_019, Version 5.3.1

Host Computer System: CompuAdd 433
under MS/DOS v5.0 and Phar Lap v4.0


Target Computer System: Intel iSBC 186/100

Customer Agreement Number: 92-10-21-ALS

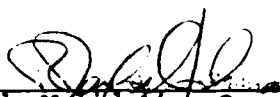
See section 3.1 for any additional information about the testing environment.

As a result of this validation effort, Validation Certificate 921210W1.11302 is awarded to Alsys, Inc. This certificate expires two years after the ANSI adoption of ANSI/MIL-STD-1815B.


This report has been reviewed and is approved.



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DECLARATION OF CONFORMANCE

Customer: Alsys Inc.

Ada Validation Facility: Wright-Patterson Air Force Base
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ACVC Version: 1.11

Ada Implementation:

Ada Compiler Name: AlsyCOMP_019

Version: 5.3.1

Host Computer System: CompuAdd 433 running MS/DOS version
5.0 and Phar Lap version 4.0.

Target Computer System: iSBC 186/100

Customer's Declaration

I, the undersigned, representing Alsys Inc, declare that Alsys Inc has no knowledge of deliberate deviations from the Ada Language Standard ANSI/MIL-STD-1815A in the implementation listed in this declaration.



Mitch Gart,
Director of Product Development

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CHAPTER 1

INTRODUCTION

The Ada implementation described above was tested according to the Ada Validation Procedures [Pro90] against the Ada Standard [Ada83] using the current Ada Compiler Validation Capability (ACVC). This Validation Summary Report (VSR) gives an account of the testing of this Ada implementation. For any technical terms used in this report, the reader is referred to [Pro90]. A detailed description of the ACVC may be found in the current ACVC User's Guide [UG89].

1.1 USE OF THIS VALIDATION SUMMARY REPORT

Consistent with the national laws of the originating country, the Ada Certification Body may make full and free public disclosure of this report. In the United States, this is provided in accordance with the "Freedom of Information Act" (5 U.S.C. #552). The results of this validation apply only to the computers, operating systems, and compiler versions identified in this report.

The organizations represented on the signature page of this report do not represent or warrant that all statements set forth in this report are accurate and complete, or that the subject implementation has no nonconformities to the Ada Standard other than those presented. Copies of this report are available to the public from the AVF which performed this validation or from:

National Technical Information Service
5285 Port Royal Road
Springfield VA 22161

Questions regarding this report or the validation test results should be directed to the AVF which performed this validation or to:

Ada Validation Organization
Computer and Software Engineering Division
Institute for Defense Analyses
1801 North Beauregard Street
Alexandria VA 22311-1772

INTRODUCTION

1.2 REFERENCES

- [Ada83] Reference Manual for the Ada Programming Language,
ANSI/MIL-STD-1815A, February 1983 and ISO 8652-1987.
- [Pro90] Ada Compiler Validation Procedures, Version 2.1, Ada Joint
Program Office, August 1990.
- [UG89] Ada Compiler Validation Capability User's Guide, 21 June 1989.

1.3 ACVC TEST CLASSES

Compliance of Ada implementations is tested by means of the ACVC. The ACVC contains a collection of test programs structured into six test classes: A, B, C, D, E, and L. The first letter of a test name identifies the class to which it belongs. Class A, C, D, and E tests are executable. Class B and class L tests are expected to produce errors at compile time and link time, respectively.

The executable tests are written in a self-checking manner and produce a PASSED, FAILED, or NOT APPLICABLE message indicating the result when they are executed. Three Ada library units, the packages REPORT and SPRT13, and the procedure CHECK FILE are used for this purpose. The package REPORT also provides a set of Identity functions used to defeat some compiler optimizations allowed by the Ada Standard that would circumvent a test objective. The package SPRT13 is used by many tests for Chapter 13 of the Ada Standard. The procedure CHECK FILE is used to check the contents of text files written by some of the Class C tests for Chapter 14 of the Ada Standard. The operation of REPORT and CHECK FILE is checked by a set of executable tests. If these units are not operating correctly, validation testing is discontinued.

Class B tests check that a compiler detects illegal language usage. Class B tests are not executable. Each test in this class is compiled and the resulting compilation listing is examined to verify that all violations of the Ada Standard are detected. Some of the class B tests contain legal Ada code which must not be flagged illegal by the compiler. This behavior is also verified.

Class L tests check that an Ada implementation correctly detects violation of the Ada Standard involving multiple, separately compiled units. Errors are expected at link time, and execution is attempted.

In some tests of the ACVC, certain macro strings have to be replaced by implementation-specific values — for example, the largest integer. A list of the values used for this implementation is provided in Appendix A. In addition to these anticipated test modifications, additional changes may be required to remove unforeseen conflicts between the tests and implementation-dependent characteristics. The modifications required for this implementation are described in section 2.3.

For each Ada implementation, a customized test suite is produced by the AVF. This customization consists of making the modifications described in the preceding paragraph, removing withdrawn tests (see section 2.1), and possibly removing some inapplicable tests (see section 2.2 and [UG89]).

In order to pass an ACVC an Ada implementation must process each test of the customized test suite according to the Ada Standard.

1.4 DEFINITION OF TERMS

Ada Compiler	The software and any needed hardware that have to be added to a given host and target computer system to allow transformation of Ada programs into executable form and execution thereof.
Ada Compiler Validation Capability (ACVC)	The means for testing compliance of Ada implementations, consisting of the test suite, the support programs, the ACVC user's guide and the template for the validation summary report.
Ada Implementation	An Ada compiler with its host computer system and its target computer system.
Ada Joint Program Office (AJPO)	The part of the certification body which provides policy and guidance for the Ada certification system.
Ada Validation Facility (AVF)	The part of the certification body which carries out the procedures required to establish the compliance of an Ada implementation.
Ada Validation Organization (AVO)	The part of the certification body that provides technical guidance for operations of the Ada certification system.
Compliance of an Ada Implementation	The ability of the implementation to pass an ACVC version.
Computer System	A functional unit, consisting of one or more computers and associated software, that uses common storage for all or part of a program and also for all or part of the data necessary for the execution of the program; executes user-written or user-designated programs; performs user-designated data manipulation, including arithmetic operations and logic operations; and that can execute programs that modify themselves during execution. A computer system may be a stand-alone unit or may consist of several inter-connected units.

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Conformity	Fulfillment by a product, process, or service of all requirements specified.
Customer	An individual or corporate entity who enters into an agreement with an AVF which specifies the terms and conditions for AVF services (of any kind) to be performed.
Declaration of Conformance	A formal statement from a customer assuring that conformity is realized or attainable on the Ada implementation for which validation status is realized.
Host Computer System	A computer system where Ada source programs are transformed into executable form.
Inapplicable test	A test that contains one or more test objectives found to be irrelevant for the given Ada implementation.
ISO	International Organization for Standardization.
LRM	The Ada standard, or Language Reference Manual, published as ANSI/MIL-STD-1815A-1983 and ISO 8652-1987. Citations from the LRM take the form "<section>.<subsection>:<paragraph>."
Operating System	Software that controls the execution of programs and that provides services such as resource allocation, scheduling, input/output control, and data management. Usually, operating systems are predominantly software, but partial or complete hardware implementations are possible.
Target Computer System	A computer system where the executable form of Ada programs are executed.
Validated Ada Compiler	The compiler of a validated Ada implementation.
Validated Ada Implementation	An Ada implementation that has been validated successfully either by AVF testing or by registration [Pro90].
Validation	The process of checking the conformity of an Ada compiler to the Ada programming language and of issuing a certificate for this implementation.
Withdrawn test	A test found to be incorrect and not used in conformity testing. A test may be incorrect because it has an invalid test objective, fails to meet its test objective, or contains erroneous or illegal use of the Ada programming language.

CHAPTER 2

IMPLEMENTATION DEPENDENCIES

2.1 WITHDRAWN TESTS

The following tests have been withdrawn by the AVO. The rationale for withdrawing each test is available from either the AVO or the AVF. The publication date for this list of withdrawn tests is 2 August 1991.

E28005C	B28006C	C32203A	C34006D	C35508I	C35508J
C35508M	C35508N	C35702A	C35702B	B41308B	C43004A
C45114A	C45346A	C45612A	C45612B	C45612C	C45651A
C46022A	B49008A	B49008B	A74006A	C74308A	B83022B
B83022H	B83025B	B83025D	C83026A	B83026B	C83041A
B85001L	C86001F	C94021A	C97116A	C98003B	BA2011A
CB7001A	CB7001B	CB7004A	CC1223A	BC1226A	CC1226B
BC3009B	BD1B02B	BD1B06A	AD1B08A	BD2A02A	CD2A21E
CD2A23E	CD2A32A	CD2A41A	CD2A41E	CD2A87A	CD2B15C
BD3006A	BD4008A	CD4022A	CD4022D	CD4024B	CD4024C
CD4024D	CD4031A	CD4051D	CD5111A	CD7004C	ED7005D
CD7005E	AD7006A	CD7006E	AD7201A	AD7201E	CD7204B
AD7206A	BD8002A	BD8004C	CD9005A	CD9005B	CDA201E
CE2107I	CE2117A	CE2117B	CE2119B	CE2205B	CE2405A
CE3111C	CE3116A	CE3118A	CE3411B	CE3412B	CE3607B
CE3607C	CE3607D	CE3812A	CE3814A	CE3902B	

2.2 INAPPLICABLE TESTS

A test is inapplicable if it contains test objectives which are irrelevant for a given Ada implementation. Reasons for a test's inapplicability may be supported by documents issued by the ISO and the AJPO known as Ada Commentaries and commonly referenced in the format AI-ddddd. For this implementation, the following tests were determined to be inapplicable for the reasons indicated; references to Ada Commentaries are included as appropriate.

IMPLEMENTATION DEPENDENCIES

The following 201 tests have floating-point type declarations requiring more digits than `SYSTEM.MAX_DIGITS`:

C24113L..Y (14 tests)	C35705L..Y (14 tests)
C35706L..Y (14 tests)	C35707L..Y (14 tests)
C35708L..Y (14 tests)	C35802L..Z (15 tests)
C45241L..Y (14 tests)	C45321L..Y (14 tests)
C45421L..Y (14 tests)	C45521L..Z (15 tests)
C45524L..Z (15 tests)	C45621L..Z (15 tests)
C45641L..Y (14 tests)	C46012L..Z (15 tests)

C35404D, C45231D, B86001X, C86006E, and CD7101G check for a predefined integer type with a name other than `INTEGER`, `LONG_INTEGER`, or `SHORT_INTEGER`; for this implementation, there is no such type.

C35713D and B86001Z check for a predefined floating-point type with a name other than `FLOAT`, `LONG_FLOAT`, or `SHORT_FLOAT`; for this implementation, there is no such type.

C45531M..P and C45532M..P (8 tests) check fixed-point operations for types that require a `SYSTEM.MAX_MANTISSA` of 47 or greater; for this implementation, `MAX_MANTISSA` is less than 47.

C45536A, C46013B, C46031B, C46033B, and C46034B contain length clauses that specify values for `'SMALL` that are not powers of two or ten; this implementation does not support such values for `'SMALL`.

C45624A..B (2 tests) check that the proper exception is raised if `MACHINE_OVERFLOW`s is `FALSE` for floating point types and the results of various floating-point operations lie outside the range of the base type; for this implementation, `MACHINE_OVERFLOW`s is `TRUE`.

B86001Y uses the name of a predefined fixed-point type other than type `DURATION`; for this implementation, there is no such type.

C96005B uses values of type `DURATION`'s base type that are outside the range of type `DURATION`; for this implementation, the ranges are the same.

CD1009C checks whether a length clause can specify a non-default size for a floating-point type; this implementation does not support such sizes.

CD2A53A checks operations of a fixed-point type for which a length clause specifies a power-of-ten `TYPE'SMALL`; this implementation does not support decimal `'SMALL`s. (See section 2.3.)

CD2A84A, CD2A84E, CD2A84I..J (2 tests), and CD2A84O use length clauses to specify non-default sizes for access types; this implementation does not support such sizes.

IMPLEMENTATION DEPENDENCIES

BD8001A, BD8003A, BD8004A..B (2 tests), and AD8011A use machine code insertions; this implementation provides no package MACHINE_CODE.

The following 264 tests check operations on sequential, text, and direct access files; this implementation does not support external files:

CE2102A..C (3)	CE2102G..H (2)	CE2102K	CE2102N..Y (12)
CE2103C..D (2)	CE2104A..D (4)	CE2105A..B (2)	CE2106A..B (2)
CE2107A..H (8)	CE2107L	CE2108A..H (8)	CE2109A..C (3)
CE2110A..D (4)	CE2111A..I (9)	CE2115A..B (2)	CE2120A..B (2)
CE2201A..C (3)	EE2201D..E (2)	CE2201F..N (9)	CE2203A
CE2204A..D (4)	CE2205A	CE2206A	CE2208B
CE2401A..C (3)	EE2401D	CE2401E..F (2)	EE2401G
CE2401H..L (5)	CE2403A	CE2404A..B (2)	CE2405B
CE2406A	CE2407A..B (2)	CE2408A..B (2)	CE2409A..B (2)
CE2410A..B (2)	CE2411A	CE3102A..C (3)	CE3102F..H (3)
CE3102J..K (2)	CE3103A	CE3104A..C (3)	CE3106A..B (2)
CE3107B	CE3108A..B (2)	CE3109A	CE3110A
CE3111A..B (2)	CE3111D..E (2)	CE3112A..D (4)	CE3114A..B (2)
CE3115A	CE3119A	EE3203A	EE3204A
CE3207A	CE3208A	CE3301A	EE3301B
CE3302A	CE3304A	CE3305A	CE3401A
CE3402A	EE3402B	CE3402C..D (2)	CE3403A..C (3)
CE3403E..F (2)	CE3404B..D (3)	CE3405A	EE3405B
CE3405C..D (2)	CE3406A..D (4)	CE3407A..C (3)	CE3408A..C (3)
CE3409A	CE3409C..E (3)	EE3409F	CE3410A
CE3410C..E (3)	EE3410F	CE3411A	CE3411C
CE3412A	EE3412C	CE3413A..C (3)	CE3414A
CE3602A..D (4)	CE3603A	CE3604A..B (2)	CE3605A..E (5)
CE3606A..B (2)	CE3704A..F (6)	CE3704M..O (3)	CE3705A..E (5)
CE3706D	CE3706F..G (2)	CE3804A..P (16)	CE3805A..B (2)
CE3806A..B (2)	CE3806D..E (2)	CE3806G..H (2)	CE3904A..B (2)
CE3905A..C (3)	CE3905L	CE3906A..C (3)	CE3906E..F (2)

CE2403A checks that WRITE raises USE_ERROR if the capacity of an external direct file is exceeded; this implementation cannot restrict file capacity.

CE3107A attempts to open a file and expect NAME_ERROR to be raised; this implementation does not support external files and so raises USE_ERROR. (See section 2.3.)

CE3304A checks that SET LINE LENGTH and SET PAGE LENGTH raise USE_ERROR if they specify an inappropriate value for the external file; there are no inappropriate values for this implementation.

CE3413B checks that PAGE raises LAYOUT_ERROR when the value of the page number exceeds COUNT'LAST; for this implementation, the value of COUNT'LAST is greater than 150000, making the checking of this objective impractical.

IMPLEMENTATION DEPENDENCIES

2.3 TEST MODIFICATIONS

Modifications (see section 1.3) were required for 28 tests.

The following tests were split into two or more tests because this implementation did not report the violations of the Ada Standard in the way expected by the original tests.

B23004A	B23007A	B23009A	B25002A	B26005A	B27005A
B28003A	B32202A..C	B37004A	B61012A	B91004A	B95069A..B
BA1101B	BC2001D	BC3009A	BC3009C	BC3205A..E	

BA2001E was graded passed by Evaluation Modification as directed by the AVO. The test expects that duplicate names of subunits with a common ancestor will be detected as compilation errors; this implementation detects the errors at link time, and the AVO ruled that this behavior is acceptable.

EA3004D was graded passed by Evaluation and Processing Modification as directed by the AVO. The test requires that either pragma INLINE is obeyed for a function call in each of three contexts and that thus three library units are made obsolete by the re-compilation of the inlined function's body, or else the pragma is ignored completely. This implementation obeys the pragma except when the call is within the package specification. When the test's files are processed in the given order, only two units are made obsolete; thus, the expected error at line 27 of file EA3004D6M is not valid and is not flagged. To confirm that indeed the pragma is not obeyed in this one case, the test was also processed with the files re-ordered so that the re-compilation follows only the package declaration (and thus the other library units will not be made obsolete, as they are compiled later); a "NOT APPLICABLE" result was produced, as expected. The revised order of files was 0-1-4-5-2-3-6.

CD2A53A was graded inapplicable by Evaluation Modification as directed by the AVO. The test contains a specification of a power-of-10 value as 'SMALL for a fixed-point type. The AVO ruled that, under ACVC 1.11, support of decimal 'SMALLs may be omitted.

CE3107A was graded inapplicable by Evaluation Modification as directed by the AVO. The test aborts with an unhandled exception when USE_ERROR is raised on the attempt to open an external file. This is acceptable behavior because this implementation does not support external files (cf. AI-00332).

CHAPTER 3
PROCESSING INFORMATION

3.1 TESTING ENVIRONMENT

The Ada implementation tested in this validation effort is described adequately by the information given in the initial pages of this report.

For technical information about this Ada implementation, contact:

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Testing of this Ada implementation was conducted at the customer's site by a validation team from the AVF.

PROCESSING INFORMATION

3.2 SUMMARY OF TEST RESULTS

An Ada Implementation passes a given ACVC version if it processes each test of the customized test suite in accordance with the Ada Programming Language Standard, whether the test is applicable or inapplicable; otherwise, the Ada Implementation fails the ACVC [Pro90].

For all processed tests (inapplicable and applicable), a result was obtained that conforms to the Ada Programming Language Standard.

The list of items below gives the number of ACVC tests in various categories. All tests were processed, except those that were withdrawn because of test errors (item b; see section 2.1), those that require a floating-point precision that exceeds the implementation's maximum precision (item e; see section 2.2), and those that depend on the support of a file system — if none is supported (item d). All tests passed, except those that are listed in sections 2.1 and 2.2 (counted in items b and f, below).

a) Total Number of Applicable Tests	3570	
b) Total Number of Withdrawn Tests	95	
c) Processed Inapplicable Tests	40	
d) Non-Processed I/O Tests	264	
e) Non-Processed Floating-Point Precision Tests	201	
f) Total Number of Inapplicable Tests	505	(c+d+e)
g) Total Number of Tests for ACVC 1.11	4170	(a+b+f)

3.3 TEST EXECUTION

Diskettes containing the customized test suite (see section 1.3) were taken on-site by the validation team for processing. The contents of the diskettes were loaded directly onto the host computer.

After the test files were loaded onto the host computer, the full set of tests was processed by the Ada implementation.

The tests were compiled and linked on the host computer system, as appropriate. The executable images were transferred to the target computer system by a RS232C communications link, and run. The results were captured on the host computer system.

Testing was performed using command scripts provided by the customer and reviewed by the validation team. See Appendix B for a complete listing of the processing options for this implementation. It also indicates the

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default options. The options invoked explicitly for validation testing during this test were:

Compiler options		Effect
CHECKS	=> ALL	Generate all execution checks.
GENERIC	=> STUBS	Do not inline generics.
TASKING	=> YES	Allow tasking.
MEMORY	=> 500	Amount of internal buffers shared by compile virtual memory.
STACK	=> 4096	Boundary size determining whether an object is allocated on the stack or in the heap.
GLOBAL	=> 4096	Boundary size determining whether an object is allocated in the global data segment or in the heap.
CALLS	=> INLINED	Allow inlining of subprograms.
REDUCTION	=> NONE	No optimization of checks or loops.
EXPRESSIONS	=> NONE	No lowlevel optimization.
Binder options		Effect
TARGET	=> CROSS	Produce cross executable
LEVEL	=> LINK	Bind and link.
EXECUTION	=> REAL	Real mode execution.
FLOAT	=> AUTOMATIC	Determine at runtime to use software or hardware floating point.
OBJECT	=> AUTOMATIC	Object name is same as main procedure (truncated to 8 characters).
UNCALLED	=> REMOVE	Remove uncalled subprograms.
SLICE	=> 0	Disable time slicing.
VERSION_ID	=> TIME_STAMP	Place time stamp in generated object.
MAIN	=> 32	Size of main program stack.
TASK	=> 4	Size of explicit Ada task stacks.
HISTORY	=> YES	Allow for stack traceback.
SIZE	=> 64	Size (in K bytes) of initial heap.
INCREMENT	=> 64	Size (in K bytes) of increment to heap.
DIRECTIVES	=> "@linkloc.rsp"	Identify linker directives file.
MODULES	=> "hw.obj"	Identify any additional object modules.
SEARCH	=> no_value	Identify any additional object libraries.

Test output, compiler and linker listings, and job logs were captured on magnetic tape and archived at the AVF. The listings examined on-site by the validation team were also archived.

APPENDIX A

MACRO PARAMETERS

This appendix contains the macro parameters used for customizing the ACVC. The meaning and purpose of these parameters are explained in [UG89]. The parameter values are presented in two tables. The first table lists the values that are defined in terms of the maximum input-line length, which is the value for \$MAX_IN_LEN—also listed here. These values are expressed here as Ada string aggregates, where "V" represents the maximum input-line length.

Macro Parameter	Macro Value
\$MAX_IN_LEN	255 — Value of V
\$BIG_ID1	(1..V-1 => 'A', V => '1')
\$BIG_ID2	(1..V-1 => 'A', V => '2')
\$BIG_ID3	(1..V/2 => 'A') & '3' & (1..V-1-V/2 => 'A')
\$BIG_ID4	(1..V/2 => 'A') & '4' & (1..V-1-V/2 => 'A')
\$BIG_INT_LIT	(1..V-3 => '0') & "298"
\$BIG_REAL_LIT	(1..V-5 => '0') & "690.0"
\$BIG_STRING1	"" & (1..V/2 => 'A') & ""
\$BIG_STRING2	"" & (1..V-1-V/2 => 'A') & '1' & ""
\$BLANKS	(1..V-20 => ' ')
\$MAX_LEN_INT_BASED_LITERAL	"2:" & (1..V-5 => '0') & "11:"
\$MAX_LEN_REAL_BASED_LITERAL	"16:" & (1..V-7 => '0') & "F.E:"

MACRO PARAMETERS

\$MAX_STRING_LITERAL ''' & (1..V-2 => 'A') & '''

The following table lists all of the other macro parameters and their respective values.

Macro Parameter	Macro Value
\$ACC_SIZE	32
\$ALIGNMENT	2
\$COUNT_LAST	2147483647
\$DEFAULT_MEM_SIZE	2**24
\$DEFAULT_STOR_UNIT	8
\$DEFAULT_SYS_NAME	I80X86
\$DELTA_DOC	2#1.0#E-31
\$ENTRY_ADDRESS	SYSTEM.VALUE("0010:0040")
\$ENTRY_ADDRESS1	SYSTEM.VALUE("0010:0080")
\$ENTRY_ADDRESS2	SYSTEM.VALUE("0010:0100")
\$FIELD_LAST	255
\$FILE_TERMINATOR	' '
\$FIXED_NAME	NO_SUCH_FIXED_NAME
\$FLOAT_NAME	NO_SUCH_FLOAT_NAME
\$FORM_STRING	""
\$FORM_STRING2	"CANNOT RESTRICT FILE CAPACITY"
\$GREATER_THAN_DURATION	75_000.0
\$GREATER_THAN_DURATION BASE LAST	I31_073.0
\$GREATER_THAN_FLOAT_BASE LAST	1.80141E+38
\$GREATER_THAN_FLOAT_SAFE LARGE	1.0E308

MACRO PARAMETERS

\$GREATER_THAN_SHORT_FLOAT_SAFE_LARGE
 1.0E308
 \$HIGH_PRIORITY 10
 \$ILLEGAL_EXTERNAL_FILE_NAME1
 \NODIRECTORY\FILENAME
 \$ILLEGAL_EXTERNAL_FILE_NAME2
 THIS_FILE_NAME_IS_TOO_LONG_FOR_MY_SYSTEM
 \$INAPPROPRIATE_LINE_LENGTH
 -1
 \$INAPPROPRIATE_PAGE_LENGTH
 -1
 \$INCLUDE_PRAGMA1 PRAGMA INCLUDE ("A28006D1.TST")
 \$INCLUDE_PRAGMA2 PRAGMA INCLUDE ("B28006D1.TST")
 \$INTEGER_FIRST -32768
 \$INTEGER_LAST 32767
 \$INTEGER_LAST_PLUS_1 32768
 \$INTERFACE_LANGUAGE C
 \$LESS_THAN_DURATION -75_000.0
 \$LESS_THAN_DURATION_BASE_FIRST
 -131_073.0
 \$LINE_TERMINATOR ASCII.CR & ASCII.LF
 \$LOW_PRIORITY 1
 \$MACHINE_CODE_STATEMENT
 NULL;
 \$MACHINE_CODE_TYPE NO_SUCH_TYPE
 \$MANTISSA_DOC 31
 \$MAX_DIGITS 15
 \$MAX_INT 2147483647
 \$MAX_INT_PLUS_1 2_147_483_648
 \$MIN_INT -2147483648
 \$NAME NO_SUCH_TYPE_AVAILABLE

MACRO PARAMETERS

\$NAME_LIST	S370,I80X86,I80386,MC680X0,VAX, TRANSPUTER,RS_6000,MIPS
\$NAME_SPECIFICATION1	D:\WORK\X2120A
\$NAME_SPECIFICATION2	D:\WORK\X2120B
\$NAME_SPECIFICATION3	D:\WORK\X3119A
\$NEG_BASED_INT	16#F000000E#
\$NEW_MEM_SIZE	2**24
\$NEW_STOR_UNIT	16
\$NEW_SYS_NAME	I80X86
\$PAGE_TERMINATOR	ASCII.CR & ASCII.LF & ASCII.FF
\$RECORD_DEFINITION	NEW INTEGER;
\$RECORD_NAME	NO_SUCH_MACHINE_CODE_TYPE
\$TASK_SIZE	32
\$TASK_STORAGE_SIZE	1024
\$TICK	1.0/18.2
\$VARIABLE_ADDRESS	SYSTEM.VALUE("0010:0020")
\$VARIABLE_ADDRESS1	SYSTEM.VALUE("0010:0024")
\$VARIABLE_ADDRESS2	SYSTEM.VALUE("0010:0028")
\$YOUR_PRAGMA	INTERFACE

APPENDIX B

COMPILATION SYSTEM OPTIONS

The compiler options of this Ada implementation, as described in this Appendix, are provided by the customer. Unless specifically noted otherwise, references in this appendix are to compiler documentation and not to this report.

```
COMPILE (SOURCE    => source_name | INSTANTIATION,
          LIBRARY   => library_name,
          OPTIONS   => (ANNOTATE  => character string,
                       ERRORS    => positive integer,
                       LEVEL     => PARSE | SEMANTIC | CODE | UPDATE,
                       CHECKS    => ALL | STACK | NONE,
                       GENERICS  => STUBS | INLINE,
                       TASKING   => YES | NO,
                       MEMORY    => number of Kbytes),
          DISPLAY   => (OUTPUT    => SCREEN | NONE | AUTOMATIC | file_name,
                       WARNING    => YES | NO,
                       TEXT       => YES | NO,
                       SHOW       => BANNER | RECAP | ALL | NONE,
                       DETAIL     => YES | NO,
                       ASSEMBLY   => CODE | MAP | ALL | NONE),
          ALLOCATION => (STACK     => positive integer,
                       GLOBAL     => positive integer),
          IMPROVE   => (CALLS     => NORMAL | INLINED,
                       REDUCTION  => NONE | PARTIAL | EXTENSIVE,
                       EXPRESSIONS=> NONE | PARTIAL | EXTENSIVE),
          KEEP      => (COPY      => YES | NO,
                       DEBUG     => YES | NO,
                       TREE      => YES | NO));
```

COMPILATION SYSTEM OPTIONS

LINKER OPTIONS

The linker options of this Ada implementation, as described in this Appendix, are provided by the customer. Unless specifically noted otherwise, references in this appendix are to linker documentation and not to this report.

```
BIND (PROGRAM    => unit_name,
      LIBRARY    => library_name,
      TARGET     => (NATIVE | CROSS | SIMULATED),
      OPTIONS    => (LEVEL      => CHECK | BIND | LINK,
                    EXECUTION => REAL | PROTECTED | EXTENDED | AUTOMATIC,
                    FLOAT     => AUTOMATIC | HARDWARE | SOFTWARE
                    OBJECT    => AUTOMATIC | file_name,
                    UNCALLED  => REMOVE | KEEP,
                    SLICE     => NO | positive_integer,
                    VERSION_ID=> TIME_STAMP | string,
                    TIMER     => NORMAL | FAST),
      STACK      => (MAIN       => positive_integer,
                    TASK       => positive_integer,
                    HISTORY    => YES | NO),
      HEAP       => (SIZE       => positive_integer,
                    INCREMENT  => positive_integer),
      INTERFACE  => (DIRECTIVES=> options_for_linker,
                    MODULES   => file_names,
                    SEARCH    => library_names),
      DISPLAY    => (OUTPUT     => SCREEN | NONE | AUTOMATIC | file_name,
                    DATA      => BIND | LINK | ALL | NONE,
                    WARNING    => YES | NO),
      KEEP       => (DEBUG      => YES | NO,
                    SYMBOLS    => NONE | PARTIAL | EXTENSIVE),
      CROSS      => (ROM        => YES | NO));
```

APPENDIX C

APPENDIX F OF THE Ada STANDARD

The only allowed implementation dependencies correspond to implementation-dependent pragmas, to certain machine-dependent conventions as mentioned in Chapter 13 of the Ada Standard, and to certain allowed restrictions on representation clauses. The implementation-dependent characteristics of this Ada implementation, as described in this Appendix, are provided by the customer. Unless specifically noted otherwise, references in this Appendix are to compiler documentation and not to this report. Implementation-specific portions of the package STANDARD, which are not a part of Appendix F, are:

package STANDARD is

```
.....
type INTEGER is range -32_768 .. 32_767;
type SHORT_INTEGER is range -128 .. 127;
type LONG_INTEGER is range -2_147_483_648 .. 2_147_483_647;

type FLOAT is digits 6 range -2#1.111_1111_1111_1111_1111_1111#E+127
.. 2#1.111_1111_1111_1111_1111_1111#E+127;
type LONG_FLOAT is digits 15 range
-2#1.1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111#1023
..
2#1.1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111#1023;
type SHORT_FLOAT is digits 6 range -2#1.111_1111_1111_1111_1111_1111#E+127
.. 2#1.111_1111_1111_1111_1111_1111#E+127;

type DURATION is delta 2#0.000_000_000_000_01#
range -131072.00000 .. 131071.99994;
```

```
.....
end STANDARD;
```

APPENDIX F OF THE Ada STANDARD

Alsys Ada Software Development Environment

for DOS to 80x86

APPENDIX F

Version 5

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APPENDIX F

Implementation - Dependent Characteristics

This appendix summarizes the implementation- dependent characteristics of the Alsys Ada Software Development Environment for DOS to 80x86. Appendix F is a required part of the Reference Manual for the Ada Programming Language (called the RM in this appendix).

The sections of this appendix are as follows:

1. The form, allowed places, and effect of every implementation-dependent pragma.
2. The name and the type of every implementation-dependent attribute.
3. The specification of the package SYSTEM.
4. The description of the representation clauses.
5. The conventions used for any implementation-generated name denoting implementation-dependent components.

6. The interpretation of expressions that appear in address clauses, including those for interrupts.
7. Any restrictions on unchecked conversions.
8. Any implementation-dependent characteristics of the input-output packages.
9. Characteristics of numeric types.
10. Other implementation-dependent characteristics.
11. Compiler limitations.

The name *Alsys Runtime Executive Programs* or simply *Runtime Executive* refers to the runtime library routines provided for all Ada programs. These routines implement the Ada heap, exceptions, tasking control, and other utility functions.

General systems programming notes are given in another document, the *Application Developer's Guide* (for example, parameter passing conventions needed for interface with assembly routines).

Section 1

1 Implementation-Dependent Pragmas

1.1 INLINE

Pragma *INLINE* is fully supported; however, it is not possible to inline a subprogram in a declarative part.

1.2 INTERFACE

Ada programs can interface with subprograms written in Assembler and other languages through the use of the predefined pragma *INTERFACE* and the implementation-defined pragma *INTERFACE_NAME*.

Pragma *INTERFACE* specifies the name of an interfaced subprogram and the name of the programming language for which parameter passing conventions will be generated. Pragma *INTERFACE* takes the form specified in the RM:

```
pragma INTERFACE (language_name, subprogram_name);
```

where,

o *language_name* is *ASSEMBLER*, *C*, or *ADA*.

APPENDIX F OF THE Ada STANDARD

- o `subprogram_name` is the name used within the Ada program to refer to the interfaced subprogram.

The only language names accepted by `pragma INTERFACE` are `ASSEMBLER`, `C`, or `Ada`. The full implementation requirements for writing `pragma INTERFACE` subprograms are described in the Application Developer's Guide.

The language name used in the `pragma INTERFACE` does not have to have any relationship to the language actually used to write the interfaced subprogram. It is used only to tell the Compiler how to generate subprogram calls; that is, what kind of parameter passing techniques to use. The programmer can interface Ada programs with subroutines written in any other (compiled) language by understanding the mechanisms used for parameter passing by the compiler and the corresponding mechanisms of the chosen external language.

1.3 `INTERFACE_NAME`

`Pragma INTERFACE NAME` associates the name of the interfaced subprogram with the external name of the interfaced subprogram. If `pragma INTERFACE NAME` is not used, then the two names are assumed to be identical. This `pragma` takes the form:

```
pragma INTERFACE_NAME (subprogram_name, string_literal);
```

where,

- o `subprogram_name` is the name used within the Ada program to refer to the interfaced subprogram.
- o `string_literal` is the name by which the interfaced subprogram is referred to at link time.

The `pragma INTERFACE NAME` is used to identify routines in other languages that are not named with legal Ada identifiers. Ada identifiers can only contain letters, digits, or underscores, whereas the DOS Linker allows external names to contain other characters, for example, the dollar sign (\$) or commercial at sign (@). These characters can be specified in the `string_literal` argument of the `pragma INTERFACE NAME`.

The `pragma INTERFACE NAME` is allowed at the same places of an Ada program as the `pragma INTERFACE`. (Location restrictions can be found in section 13.9 of the RM.) However, the `pragma INTERFACE NAME` must always occur after the `pragma INTERFACE` declaration for the interfaced subprogram.

The string literal of the `pragma INTERFACE NAME` is passed through unchanged to the DOS object file. The maximum length of the `string_literal` is 40 characters. This limit is not checked by the Compiler, but the string is truncated by the Binder to meet the Intel object module format standard.

The user must be aware however, that some tools from other vendors do not fully support the standard object file format and may restrict the length of symbols.

For example, the IBM and Microsoft assemblers silently truncate symbols at 31 characters.

The Runtime Executive contains several external identifiers. All such identifiers begin with either the string "ADA@" or the string "ADAS@". Accordingly, names prefixed by "ADA@" or "ADAS@" should be avoided by the user.

Example

```
package SAMPLE_DATA is
  function SAMPLE_DEVICE (X: INTEGER) return
  INTEGER;
  function PROCESS_SAMPLE (X: INTEGER) return
  INTEGER;
private
  pragma INTERFACE (ASSEMBLER, SAMPLE_DEVICE);
  pragma INTERFACE (ADA, PROCESS_SAMPLE);
  pragma INTERFACE_NAME (SAMPLE_DEVICE, "DEVIO$GET_SAMPLE");
end SAMPLE_DATA;
```

1.4 INDENT

Pragma INDENT is only used with AdaReformat. AdaReformat is the Alsys reformatter which offers the functionalities of a pretty-printer in an Ada environment.

The pragma is placed in the source file and interpreted by the Reformatter. The line

```
pragma INDENT(OFF);
```

causes AdaReformat not to modify the source lines after this pragma, while

```
pragma INDENT(ON);
```

causes AdaReformat to resume its action after this pragma.

1.5 Other Pragas

Pragmas IMPROVE and PACK are discussed in detail in the section on representation clauses and records (Chapter 4).

Pragma PRIORITY is accepted with the range of priorities running from 1 to 10 (see the definition of the predefined package SYSTEM in Section 3). Undefined priority (no pragma PRIORITY) is treated as though it were less than any defined priority value.

In addition to pragma SUPPRESS, it is possible to suppress all checks in a given compilation by the use of the Compiler option CHECKS. (See Chapter 4 of the User's Guide.)

2 Implementation-Dependent Attributes

APPENDIX F OF THE Ada STANDARD

2.1 P'IS_ARRAY

For a prefix P that denotes any type or subtype, this attribute yields the value TRUE if P is an array type or an array subtype; otherwise, it yields the value FALSE.

2.2 E'EXCEPTION_CODE

For a prefix E that denotes an exception name, this attribute yields a value that represents the internal code of the exception. The value of this attribute is of the type INTEGER.

2.3 Attributes Used in Record Representation Clauses

In addition to the Representation Attributes of [13.7.2] and [13.7.3], the following attributes are used to form names of indirect and implicit components for use in record representation clauses, as described in Section 4.8.

- 'OFFSET
- 'RECORD_SIZE
- 'VARIANT_INDEX
- 'ARRAY_DESCRIPTOR
- 'RECORD_DESCRIPTION

3 Specification of the package SYSTEM

This file is provided for documentation purposes only. It is already compiled into the predefined library. Do not re-compile it.

```
-----
*****
--
--  COPYRIGHT NOTICE:
--  "This unpublished work is protected both as a proprietary work and
--  under the Universal Copyright Convention and the US Copyright Act of
--  1976. Its distribution and access are limited only to authorized
--  persons. Copr. (c) Alsys. Created 1990, initially licensed 1990.
--  All rights reserved.
--
--  Unauthorized use (including use to prepare other works), disclosure,
--  reproduction, or distribution may violate national criminal law."
--
*****
```

package SYSTEM is

— The order of the elements of this type is not significant.

type NAME is (I80X86);

SYSTEM_NAME : constant NAME := I80X86;

STORAGE_UNIT : constant := 8;

MAX_INT : constant := $2^{31} - 1$;

MIN_INT : constant := $-(2^{31})$;

MAX_MANTISSA : constant := 31;

FINE_DELTA : constant := $2 \times 1.0 \times 10^{-31}$;

MAX_DIGITS : constant := 15;

MEMORY_SIZE : constant := 2^{24} ;

TICK : constant := $1.0 / 18.2$;

subtype PRIORITY is INTEGER range 1 .. 10;

type ADDRESS is private;

NULL_ADDRESS : constant ADDRESS;

— Converts a string to an address. The syntax of the string and its meaning are target dependent.

—

— For the 8086, 80186 and 80286 the syntax is:

— "SSSS:OOOO" where SSSS and OOOO are a 4 digit or less hexadecimal number representing a segment value and an offset.

— The physical address corresponding to SSSS:OOOO depends on the execution mode. In real mode it is $16 \times \text{SSSS} + \text{OOOO}$.

— In protected mode the value SSSS represents a segment descriptor.

— Example:

— "0014:00F0"

—

— The exception CONSTRAINT_ERROR is raised if the string has not the proper syntax.

function VALUE (LEFT : in STRING) return ADDRESS;

— Converts an address to a string. The syntax of the returned string is described in the VALUE function.

APPENDIX F OF THE Ada STANDARD

```
subtype ADDRESS_STRING is STRING(1..9);
```

```
function IMAGE (LEFT : in ADDRESS) return ADDRESS_STRING;
```

- The following routines provide support to perform address
- computation. The meaning of the "+" and "-" operators is
- architecture dependent. For example on a segmented machine
- the OFFSET parameter is added to, or subtracted from the offset
- part of the address, the segment remaining untouched.

```
type OFFSET is range 0 .. 2**16 -1;
```

- On a segmented architecture the function returns true if the
- two addresses have the same segment value. On a non segmented
- architecture it always returns TRUE.
- Only makes sense for 8086, 80186, and 80286.

```
function SAME_SEGMENT (LEFT, RIGHT : in ADDRESS) return BOOLEAN;
```

- The exception ADDRESS_ERROR is raised by "<", "<=", ">", ">=", "-"
- if the two addresses do not have the same segment value. This
- exception is never raised on a non segmented machine.
- The exception CONSTRAINT_ERROR can be raised by "+" and "-".

```
ADDRESS_ERROR : exception;
```

```
function "+" (LEFT : in ADDRESS; RIGHT : in OFFSET) return ADDRESS;
function "+" (LEFT : in OFFSET; RIGHT : in ADDRESS) return ADDRESS;
function "-" (LEFT : in ADDRESS; RIGHT : in OFFSET) return ADDRESS;
```

- The exception ADDRESS_ERROR is raised on a segmented architecture
- if the two addresses do not have the same segment value.

```
function "-" (LEFT : in ADDRESS; RIGHT : in ADDRESS) return OFFSET;
```

- Perform an unsigned comparison on addresses or offset part of
- addresses on a segmented machine.

```
function "<=" (LEFT, RIGHT : in ADDRESS) return BOOLEAN;
function "<" (LEFT, RIGHT : in ADDRESS) return BOOLEAN;
function ">=" (LEFT, RIGHT : in ADDRESS) return BOOLEAN;
function ">" (LEFT, RIGHT : in ADDRESS) return BOOLEAN;
function "mod" (LEFT : in ADDRESS; RIGHT : in POSITIVE) return NATURAL;
```

- Returns the given address rounded to a specific value.

```
type ROUND_DIRECTION is (DOWN, UP);
```

```
function ROUND (VALUE : in ADDRESS;
DIRECTION : in ROUND_DIRECTION;
MODULUS : in POSITIVE) return ADDRESS;
```

- These routines are provided to perform READ/WRITE operation
- in memory.
- Warning: These routines will give unexpected results if used with
- unconstrained types.

```
generic
  type TARGET is private;
  function FETCH_FROM_ADDRESS (A : in ADDRESS)
return TARGET;
```

```
generic
  type TARGET is private;
  procedure ASSIGN_TO_ADDRESS (A : in ADDRESS; T : in TARGET);
```

- Procedure to copy LENGTH storage unit starting at the address
- FROM to the address TO. The source and destination may overlap.
- OBJECT_LENGTH designates the size of an object in storage units.

```
type OBJECT_LENGTH is range 0 .. 2**16 -1;
for OBJECT_LENGTH'SIZE use 16;
```

```
procedure MOVE (TO      : in ADDRESS;
  FROM      : in ADDRESS;
  LENGTH    : in OBJECT_LENGTH);
```

end SYSTEM;

4 Support for Representation Clauses

This section explains how objects are represented and allocated and how it is possible to control this using representation clauses. Applicable restrictions on representation clauses are also described.

The representation of an object is closely connected with its type. For this reason this section addresses successively the representation of enumeration, integer, floating point, fixed point, access, task, array and record types. For each class of type the representation of the corresponding objects is described.

Except in the case of array and record types, the description for each class of type is independent of the others. To understand the representation of array and record types it is necessary to understand first the representation of their components.

Apart from implementation defined pragmas, Ada provides three means to control the size of objects:

- o a (predefined) pragma PACK, applicable to array types

APPENDIX F OF THE Ada STANDARD

- o a record representation clause
- o a size specification

For each class of types the effect of a size specification is described. Interactions among size specifications, packing and record representation clauses is described under the discussion of array and record types.

Representation clauses on derived record types or derived tasks types are not supported.

Size representation clauses on types derived from private types are not supported when the derived type is declared outside the private part of the defining package.

4.1 Enumeration Types

4.1.1 Enumeration Literal Encoding

When no enumeration representation clause applies to an enumeration type, the internal code associated with an enumeration literal is the position number of the enumeration literal. Then, for an enumeration type with n elements, the internal codes are the integers $0, 1, 2, \dots, n-1$.

An enumeration representation clause can be provided to specify the value of each internal code as described in RM 13.3. The Alsys compiler fully implements enumeration representation clauses.

As internal codes must be machine integers the internal codes provided by an enumeration representation clause must be in the range $-(2^{31}) \dots (2^{31})-1$.

An enumeration value is always represented by its internal code in the program generated by the compiler.

4.1.2 Enumeration Types and Object Sizes

Minimum size of an enumeration subtype

The minimum possible size of an enumeration subtype is the minimum number of bits that is necessary for representing the internal codes of the subtype values in normal binary form.

A static subtype, with a null range has a minimum size of 1. Otherwise, if m and M are the values of the internal codes associated with the first and last enumeration values of the subtype, then its minimum size L is determined as follows. For $m \geq 0$, L is the smallest positive integer such that $M \leq 2^{L-1}$. For $m < 0$, L is the smallest positive integer such that $-2^{L-1} \leq m$ and $M \leq 2^{L-1}-1$. For example:

```
type COLOR is (GREEN, BLACK, WHITE, RED, BLUE, YELLOW);  
— The minimum size of COLOR is 3 bits.
```



```

subtype BLACK AND WHITE is COLOR range BLACK .. WHITE;
-- The minimum size of BLACK AND WHITE is 2 bits.
subtype BLACK OR WHITE is BLACK AND WHITE range X .. X;
-- Assuming that X is not static, the minimum
size of BLACK OR WHITE is
-- 2 bits (the same as the minimum size of its
type mark BLACK_AND_WHITE).

```

Size of an enumeration subtype

When no size specification is applied to an enumeration type or first named subtype, the objects of that type or first named subtype are represented as signed machine integers. The machine provides 8, 16 and 32 bit integers, and the compiler selects automatically the smallest signed machine integer which can hold each of the internal codes of the enumeration type (or subtype). The size of the enumeration type and of any of its subtypes is thus 8, 16 or 32 bits.

When a size specification is applied to an enumeration type, this enumeration type and each of its subtypes has the size specified by the length clause. The same rule applies to a first named subtype. The size specification must of course specify a value greater than or equal to the minimum size of the type or subtype to which it applies:

```

type EXTENDED is
  (— The usual ASCII character set.
   NUL, SOH, STX, ETX, EOT, ENQ, ACK, BEL,
   ...
   'x', 'y', 'z', '{', '|', '}', '~', DEL,
   — Extended characters
   C_CEDILLA_CAP, U_UMLAUT, E_ACUTE, ...);

```

for EXTENDED'SIZE use 8;

```

-- The size of type EXTENDED will be one byte.  Its objects will be represented
-- as unsigned 8 bit integers.

```

The Alsys compiler fully implements size specifications. Nevertheless, as enumeration values are coded using integers, the specified length cannot be greater than 32 bits.

Size of the objects of an enumeration subtype

Provided its size is not constrained by a record component clause or a pragma PACK, an object of an enumeration subtype has the same size as its subtype.

4.2 Integer Types

APPENDIX F OF THE Ada STANDARD

There are three predefined integer types in the Alsys Software Development Environment for DOS to 80x86 implementation:

```
type SHORT_INTEGER    is range -2**07 .. 2**07-1;
type INTEGER          is range -2**15 .. 2**15-1;
type LONG_INTEGER     is range -2**31 .. 2**31-1;
```

4.2.1 Integer Type Representation

An integer type declared by a declaration of the form:

```
type T is range L .. R;
```

is implicitly derived from a predefined integer type. The compiler automatically selects the predefined integer type whose range is the smallest that contains the values L to R inclusive. Binary code is used to represent integer values. Negative numbers are represented using two's complement.

4.2.2 Integer Type and Object Size

Minimum size of an integer subtype

The minimum possible size of an integer subtype is the minimum number of bits that is necessary for representing the internal codes of the subtype values in normal binary form.

For a static subtype, if it has a null range its minimum size is 1. Otherwise, if m and M are the lower and upper bounds of the subtype, then its minimum size L is determined as follows. For $m \geq 0$, L is the smallest positive integer such that $M \leq 2^{(L-1)}$. For $m < 0$, L is the smallest positive integer that $-2^{(L-1)} \leq m$ and $M \leq 2^{(L-1)} - 1$. For example:

```
subtype S is INTEGER range 0 .. 7;
— The minimum size of S is 3 bits.
subtype D is S range X .. Y;    — Assuming that X and Y are not static, the
minimum size of
— D is 3 bits (the same as the minimum size of
its type mark S).
```

Size of an integer subtype

The sizes of the predefined integer types SHORT_INTEGER, INTEGER and LONG_INTEGER are respectively 8, 16 and 32 bits.

When no size specification is applied to an integer type or to its first named subtype (if any), its size and the size of any of its subtypes is the size of the predefined type from which it derives, directly or indirectly. For example:

```
type S is range 80 .. 100;
— S is derived from SHORT_INTEGER, its size is 8
```

bits.

```
type J is range 0 .. 255;
— J is derived from INTEGER, its size is 16
bits.
```

```
type N is new J range 80 .. 100;
— N is indirectly derived from INTEGER, its size
is 16 bits.
```

When a size specification is applied to an integer type, this integer type and each of its subtypes has the size specified by the length clause. The same rule applies to a first named subtype. The size specification must of course specify a value greater than or equal to the minimum size of the type or subtype to which it applies:

```
type S is range 80 .. 100; for S'SIZE use 32;
— S is derived from SHORT_INTEGER, but its size
is 32 bits
— because of the size specification.
```

```
type J is range 0 .. 255;
for J'SIZE use 8;
— J is derived from INTEGER, but its size is 8
bits because
— of the size specification.
```

```
type N is new J range 80 .. 100;
— N is indirectly derived from INTEGER, but its
size is
— 8 bits because N inherits the size
specification of J.
```

Size of the objects of an integer subtype

Provided its size is not constrained by a record component clause or a pragma PACK, an object of an integer subtype has the same size as its subtype.

4.3 Floating Point Types

There are three predefined floating point types in the Alsys implementation for I80x86 machines:

```
type SHORT FLOAT is
  digits 7 range  $-(2.0 - 2.0^{**}(-23)) \cdot 2.0^{**127} ..$ 
   $(2.0 - 2.0^{**}(-23)) \cdot 2.0^{**127};$ 
```

```
type FLOAT is
  digits 6 range  $-(2.0 - 2.0^{**}(-23)) \cdot 2.0^{**127} ..$ 
   $(2.0 - 2.0^{**}(-23)) \cdot 2.0^{**127};$ 
```

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```
type LONG_FLOAT is
  digits_15 range -(2.0 - 2.0**(-51))*2.0**1023 ..
    (2.0 - 2.0**(-51))*2.0**1023;
```

Note that SHORT_FLOAT has the same range as FLOAT.

4.3.1 Floating Point Type Representation

A floating point type declared by a declaration of the form:

```
type T is digits D [range L .. R];
```

is implicitly derived from a predefined floating point type. The compiler automatically selects the smallest predefined floating point type whose number of digits is greater than or equal to D and which contains the values L to R inclusive.

In the program generated by the compiler, floating point values are represented using the IEEE standard formats for single and double floats.

The values of the predefined types SHORT_FLOAT and FLOAT are represented using the single float format. The values of the predefined type LONG_FLOAT are represented using the double float format. The values of any other floating point type are represented in the same way as the values of the predefined type from which it derives, directly or indirectly.

4.3.2 Floating Point Type and Object Size

The minimum possible size of a floating point subtype is 32 bits if its base type is SHORT_FLOAT or FLOAT or a type derived from SHORT_FLOAT or FLOAT; it is 64 bits if its base type is LONG_FLOAT or a type derived from LONG_FLOAT.

The sizes of the predefined floating point types SHORT_FLOAT and FLOAT is 32 bits and LONG_FLOAT is 64 bits.

The size of a floating point type and the size of any of its subtypes is the size of the predefined type from which it derives directly or indirectly.

The only size that can be specified for a floating point type or first named subtype using a size specification is its usual size (32 or 64 bits).

An object of a floating point subtype has the same size as its subtype.

4.4 Fixed Point Types

4.4.1 Fixed Point Type Representation

If no specification of small applies to a fixed point type, then the value of small is determined by the value of delta as defined by RM 3.5.9.

A specification of small can be used to impose a value of small. The value of small is required to be a power of two.

To implement fixed point types, the compiler uses a set of anonymous predefined types of the form:

```
type SHORT_FIXED is delta D range  $(-2.0^{**7}-1)*S$  ..  $2.0^{**7}*S$ ;
for SHORT_FIXED'SMALL use S;
```

```
type FIXED is delta D range  $(-2.0^{**15}-1)*S$  ..  $2.0^{**15}*S$ ;
for FIXED'SMALL use S;
```

```
type LONG_FIXED is delta D range  $(-2.0^{**31}-1)*S$  ..  $2.0^{**31}*S$ ;
for LONG_FIXED'SMALL use S;
```

where D is any real value and S any power of two less than or equal to D.

A fixed point type declared by a declaration of the form:

```
type T is delta D range L .. R;
```

possibly with a small specification:

```
for T'SMALL use S;
```

is implicitly derived from a predefined fixed point type. The compiler automatically selects the predefined fixed point type whose small and delta are the same as the small and delta of T and whose range is the shortest that includes the values L to R inclusive.

In the program generated by the compiler, a safe value V of a fixed point subtype F is represented as the integer:

```
V / F'BASE'SMALL
```

4.4.2 Fixed Point Type and Object Size

The minimum possible size of a fixed point subtype is the minimum number of binary digits that is necessary for representing the values of the range of the subtype using the small of the base type.

For a static subtype, if it has a null range its minimum size is 1. Otherwise, s and S being the bounds of the subtype, if i and I are the integer representations of m and M, the smallest and the greatest model numbers of the base type such that $s < m$ and $M < S$, then the minimum size L is determined as follows. For $i \geq 0$, L is the smallest positive integer such that $I \leq 2^{*(L-1)}$. For $i < 0$, L is the smallest positive integer such that $-2^{*(L-1)} \leq i$ and $I \leq 2^{*(L-1)}-1$.

```
type F is delta 2.0 range 0.0 .. 500.0;
— The minimum size of F is 8 bits.
```

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subtype S is F delta 16.0 range 0.0 .. 250.0;
— The minimum size of S is 7 bits.

subtype D is S range X .. Y;
— Assuming that X and Y are not static, the minimum size of D is 7 bits
— (the same as the minimum size of its type mark S).

Size of a fixed point subtype

The sizes of the predefined fixed point types `SHORT_FIXED`, `FIXED` and `LONG_FIXED` are respectively 8, 16 and 32 bits.

When no size specification is applied to a fixed point type or to its first named subtype, its size and the size of any of its subtypes is the size of the predefined type from which it derives directly or indirectly. For example:

type S is delta 0.01 range 0.8 .. 1.0;
— S is derived from an 8 bit predefined fixed type, its size is 8 bits.

type F is delta 0.01 range 0.0 .. 2.0;
— F is derived from a 16 bit predefined fixed type, its size is 16 bits.

type N is new F range 0.8 .. 1.0;
— N is indirectly derived from a 16 bit predefined fixed type,
— its size is 16 bits.

When a size specification is applied to a fixed point type, this fixed point type and each of its subtypes has the size specified by the length clause. The same rule applies to a first named subtype. The size specification must of course specify a value greater than or equal to the minimum size of the type or subtype to which it applies:

type S is delta 0.01 range 0.8 .. 1.0;
for S'SIZE use 32;
— S is derived from an 8 bit predefined fixed type, but its size is 32 bits
— because of the size specification.

type F is delta 0.01 range 0.0 .. 2.0;
for F'SIZE use 8;
— F is derived from a 16 bit predefined fixed type, but its size is 8 bits
— because of the size specification.

type N is new F range 0.8 .. 1.0;
— N is indirectly derived from a 16 bit predefined fixed type,
— but its size is 8 bits because N inherits the size specification of F.

The Alsys compiler fully implements size specifications. Nevertheless, as fixed point objects are represented using machine integers, the specified length cannot be greater than 32 bits.

Provided its size is not constrained by a record component clause or a pragma `PACK`, an object of a fixed point type has the same size as its subtype.

4.5 Access Types and Collections

Access Types and Objects of Access Types

The only size that can be specified for an access type using a size specification is its usual size (32 bits).

An object of an access subtype has the same size as its subtype, thus an object of an access subtype is always 32 bits long.

Collection Size

As described in RM 13.2, a specification of collection size can be provided in order to reserve storage space for the collection of an access type.

When no `STORAGE_SIZE` specification applies to an access type, no storage space is reserved for its collection, and the value of the attribute `STORAGE_SIZE` is then 0.

The maximum size allowed for a collection is 64kbytes.

4.6 Task Types

Storage for a task activation

As described in RM 13.2, a length clause can be used to specify the storage space (that is, the stack size) for the activation of each of the tasks of a given type. `Alsys` also allows the task stack size, for all tasks, to be established using a Binder option. If a length clause is given for a task type, the value indicated at bind time is ignored for this task type, and the length clause is obeyed. When no length clause is used to specify the storage space to be reserved for a task activation, the storage space indicated at bind time is used for this activation.

A length clause may not be applied to a derived task type. The same storage space is reserved for the activation of a task of a derived type as for the activation of a task of the parent type.

The minimum size of a task subtype is 32 bits.

A size specification has no effect on a task type. The only size that can be specified using such a length clause is its usual size (32 bits).

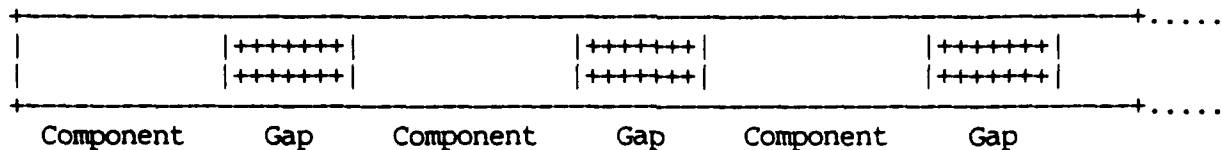
An object of a task subtype has the same size as its subtype. Thus an object of a task subtype is always 32 bits long.

4.7 Array Types

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Each array is allocated in a contiguous area of storage units. All the components have the same size. A gap may exist between two consecutive components (and after the last one). All the gaps have the same size.

4.7.1 Array Layout and Structure and Pragma PACK



If pragma PACK is not specified for an array, the size of the components is the size of the subtype of the components:

```
type A is array (1 .. 8) of BOOLEAN;
-- The size of the components of A is the size of the type BOOLEAN: 8 bits.

type DECIMAL DIGIT is range 0 .. 9;
for DECIMAL DIGIT'SIZE use 4;
type BINARY CODED DECIMAL is
  array (INTEGER range <>) of DECIMAL DIGIT;
-- The size of the type DECIMAL DIGIT is 4 bits. Thus in an array of
-- type BINARY CODED DECIMAL each component will be represented on
-- 4 bits as in the usual BCD representation.
```

If pragma PACK is specified for an array and its components are neither records nor arrays, the size of the components is the minimum size of the subtype of the components:

```
type A is array (1 .. 8) of BOOLEAN;
pragma PACK(A);
-- The size of the components of A is the minimum size of the type BOOLEAN:
-- 1 bit.

type DECIMAL DIGIT is range 0 .. 9;
for DECIMAL DIGIT'SIZE use 32;
type BINARY CODED DECIMAL is
  array (INTEGER range <>) of DECIMAL DIGIT;
pragma PACK(BINARY CODED DECIMAL);
-- The size of the type DECIMAL DIGIT is 32 bits, but, as
-- BINARY CODED DECIMAL is packed, each component of an array of this
-- type will be represented on 4 bits as in the usual BCD representation.
```

Packing the array has no effect on the size of the components when the components are records or arrays, since records and arrays may be assigned addresses consistent with the alignment of their subtypes.

Gaps

If the components are records or arrays, no size specification applies to the subtype of the components and the array is not packed, then the compiler may choose a representation with a gap after each component; the aim of the insertion of such gaps is to optimize access to the array components and to their subcomponents. The size of the gap is chosen so that the relative displacement of consecutive components is a multiple of the alignment of the subtype of the components. This strategy allows each component and subcomponent to have an address consistent with the alignment of its subtype:

```

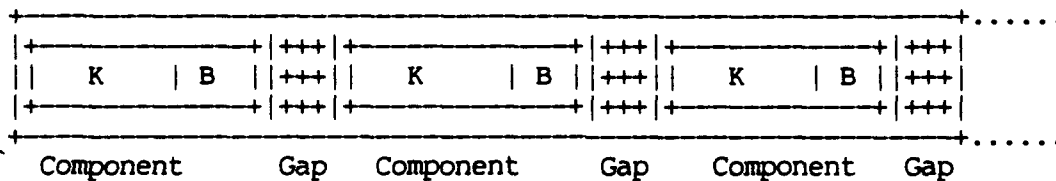
type R is
  record
    K : INTEGER;
    B : BOOLEAN;
  end record;
for R use
  record
    K at 0 range 0 .. 15;
    B at 2 range 0 .. 0;
  end record;
-- Record type R is byte aligned. Its size is 17 bits.

```

```

type A is array (1 .. 10) of R;
-- A gap of 7 bits is inserted after each component in order to respect the
-- alignment of type R. The size of an array of type A will be 240 bits.

```



Array of type A: each subcomponent K has an even offset.

If a size specification applies to the subtype of the components or if the array is packed, no gaps are inserted:

```

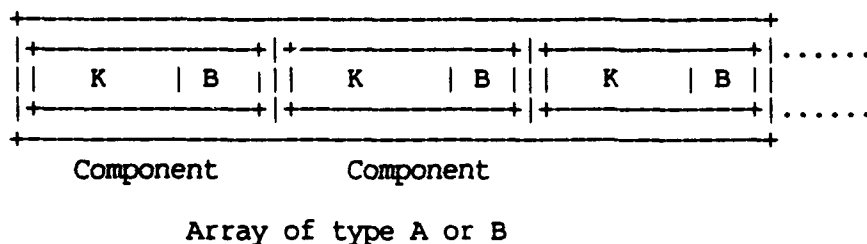
type R is
  record
    K : INTEGER;
    B : BOOLEAN;
  end record;

type A is array (1 .. 10) of R;
pragma PACK(A);
-- There is no gap in an array of type A because A is packed.
-- The size of an object of type A will be 240 bits.

type NR is new R;   for NR'SIZE use 24;

```

- type B is array (1 .. 10) of NR;
- There is no gap in an array of type B because
 - NR has a size specification.
 - The size of an object of type B will be 240 bits.



4.7.2 Array Subtype and Object Size

Size of an array subtype

The size of an array subtype is obtained by multiplying the number of its components by the sum of the size of the components and the size of the gaps (if any). If the subtype is unconstrained, the maximum number of components is considered.

The size of an array subtype cannot be computed at compile time

- o if it has non-static constraints or is an unconstrained array type with non-static index subtypes (because the number of components can then only be determined at run time).
- o if the components are records or arrays and their constraints or the constraints of their subcomponents (if any) are not static (because the size of the components and the size of the gaps can then only be determined at run time).

As has been indicated above, the effect of a pragma PACK on an array type is to suppress the gaps. The consequence of packing an array type is thus to reduce its size.

If the components of an array are records or arrays and their constraints or the constraints of their subcomponents (if any) are not static, the compiler ignores any pragma PACK applied to the array type but issues a warning message. Apart from this limitation, array packing is fully implemented by the Alsyp compiler.

A size specification applied to an array type or first named subtype has no effect. The only size that can be specified using such a length clause is its usual size. Nevertheless, such a length clause can be useful to verify that the

layout of an array is as expected by the application.

Size of the objects of an array subtype

The size of an object of an array subtype is always equal to the size of the subtype of the object.

4.8 Record Types

4.8.1 Basic Record Structure

Layout of a record

Each record is allocated in a contiguous area of storage units. The size of a record component depends on its type.

The positions and the sizes of the components of a record type object can be controlled using a record representation clause as described in RM 13.4. In the implementation there is no restriction on the position that can be specified for a component of a record. If a component is not a record or an array, its size can be any size from the minimum size to the size of its subtype. If a component is a record or an array, its size must be the size of its subtype.

```

type DEVICE_INFO_RECORD is
  record
    BIT15  : BOOLEAN;    -- Bit 15 (reserved)
    CTRL   : BOOLEAN;    -- Bit 14 (true if control strings processed)
    NETWORK : BOOLEAN;    -- Bit 13 (true if device is on network)
    BIT12  : BOOLEAN;    -- Bit 12 (reserved)
    BIT11  : BOOLEAN;    -- Bit 11 (reserved)
    BIT10  : BOOLEAN;    -- Bit 10 (reserved)
    BIT9   : BOOLEAN;    -- Bit 9 (reserved)
    BIT8   : BOOLEAN;    -- Bit 8 (reserved)
    ISDEV  : BOOLEAN;    -- Bit 7 (true if device, false if disk file)
    EOF    : BOOLEAN;    -- Bit 6 (true if at end of file)
    BINARY : BOOLEAN;    -- Bit 5 (true if binary (raw) mode)
    BIT4   : BOOLEAN;    -- Bit 4 (reserved)
    ISCLK  : BOOLEAN;    -- Bit 3 (true if clock device)
    ISNUL  : BOOLEAN;    -- Bit 2 (true if NUL device)
    ISCOT  : BOOLEAN;    -- Bit 1 (true if console output device)
    ISCIN  : BOOLEAN;    -- Bit 0 (true if console input device)
  end record;

```

```

for DEVICE_INFO_RECORD use
  record
    BIT15  at 1 range 7 .. 7;  -- Bit 15
    CTRL   at 1 range 6 .. 6;  -- Bit 14
    NETWORK at 1 range 5 .. 5;  -- Bit 13
    BIT12  at 1 range 4 .. 4;  -- Bit 12
    BIT11  at 1 range 3 .. 3;  -- Bit 11
    BIT10  at 1 range 2 .. 2;  -- Bit 10

```

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```

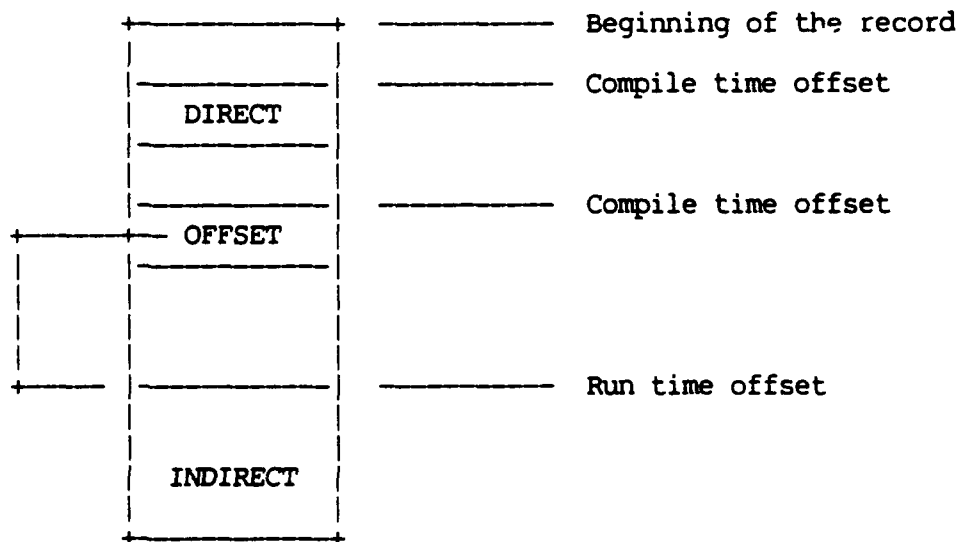
BIT9    at 1 range 1 .. 1;  — Bit 9
BIT8    at 1 range 0 .. 0;  — Bit 8
ISDEV   at 0 range 7 .. 7;  — Bit 7
EOF     at 0 range 6 .. 6;  — Bit 6
BINARY  at 0 range 5 .. 5;  — Bit 5
BIT4    at 0 range 4 .. 4;  — Bit 4
ISCLK   at 0 range 3 .. 3;  — Bit 3
ISNUL   at 0 range 2 .. 2;  — Bit 2
ISCOT   at 0 range 1 .. 1;  — Bit 1
ISCIN   at 0 range 0 .. 0;  — Bit 0
end record;
```

Pragma PACK has no effect on records. It is unnecessary because record representation clauses provide full control over record layout.

A record representation clause need not specify the position and the size for every component. If no component clause applies to a component of a record, its size is the size of its subtype.

4.8.2 Indirect Components 'OFFSET

If the offset of a component cannot be computed at compile time, this offset is stored in the record objects at run time and used to access the component. Such a component is said to be indirect while other components are said to be direct:



A direct and an indirect component

If a record component is a record or an array, the size of its subtype may be evaluated at run time and may even depend on the discriminants of the record. We will call these components dynamic components:

```
type DEVICE is (SCREEN, PRINTER);
```

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```

type COLOR is (GREEN, RED, BLUE);

type SERIES is array (POSITIVE range <>) of
INTEGER;

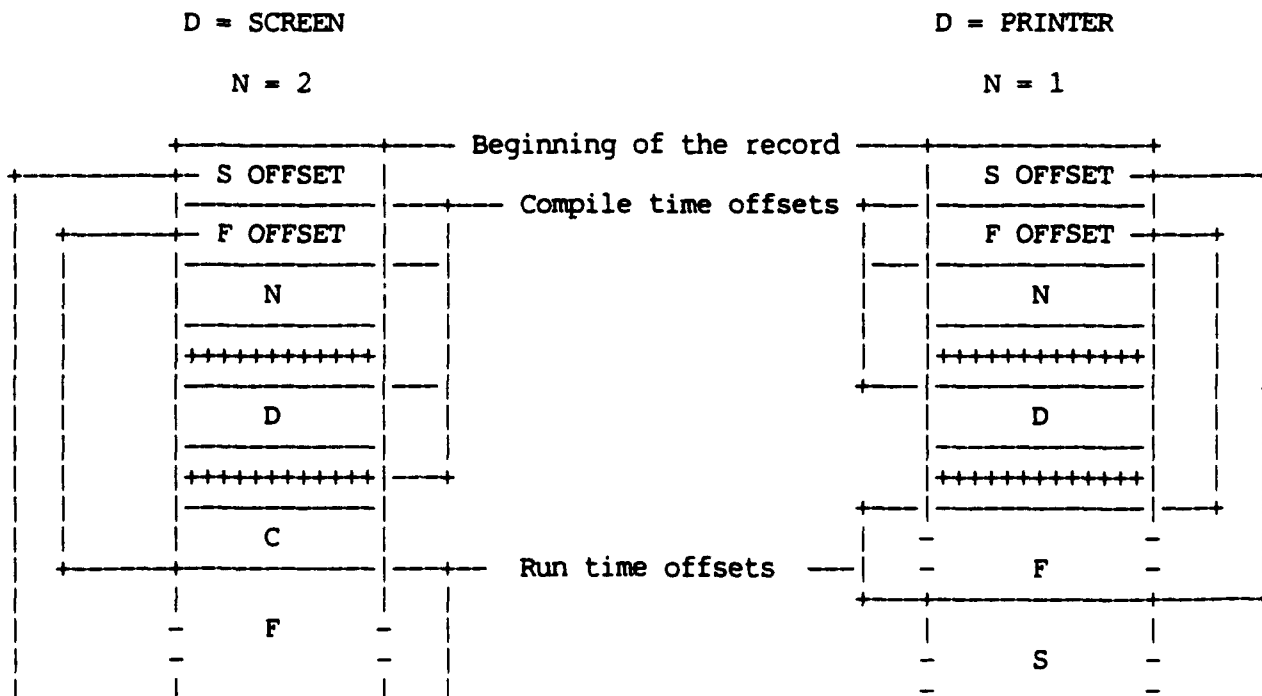
type GRAPH (L : NATURAL) is
  record
    X : SERIES(1 .. L); -- The size of X depends on L
    Y : SERIES(1 .. L); -- The size of Y depends on L
  end record;

Q : POSITIVE;

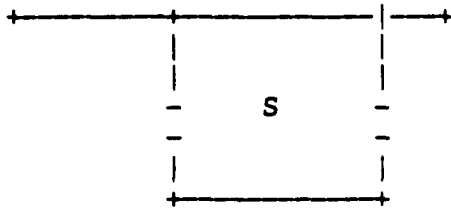
type PICTURE (N : NATURAL; D : DEVICE) is
  record
    F : GRAPH(N); -- The size of F depends on N
    S : GRAPH(Q); -- The size of S depends on Q
    case D is
      when SCREEN =>
        C : COLOR;
      when PRINTER =>
        null;
    end case;
  end record;

```

Any component placed after a dynamic component has an offset which cannot be evaluated at compile time and is thus indirect. In order to minimize the number of indirect components, the compiler groups the dynamic components together and places them at the end of the record:



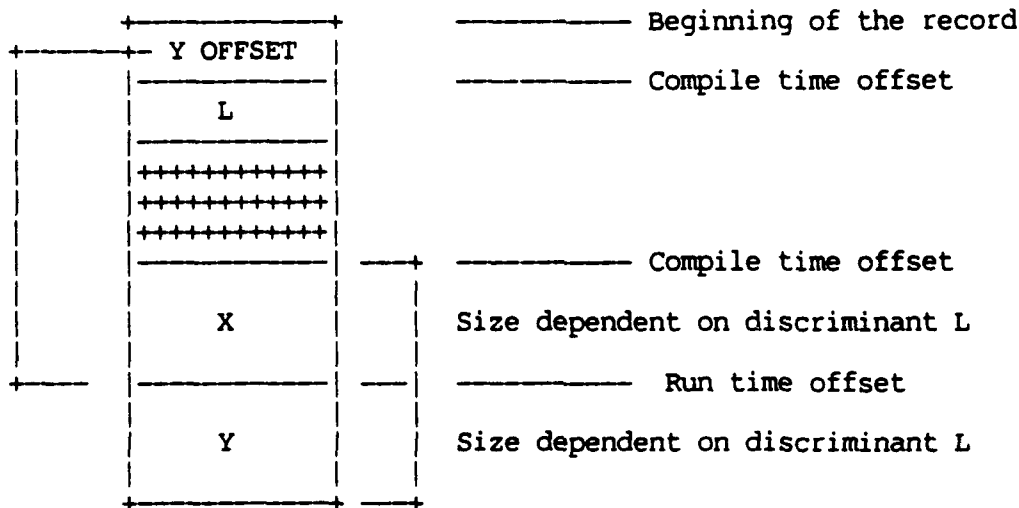
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The record type PICTURE: F and S are placed at the end of the record

Note that Ada does not allow representation clauses for record components with non-static bounds [RM 13.4.7], so the compiler's grouping of dynamic components does not conflict with the use of representation clauses.

Because of this approach, the only indirect components are dynamic components. But not all dynamic components are necessarily indirect: if there are dynamic components in a component list which is not followed by a variant part, then exactly one dynamic component of this list is a direct component because its offset can be computed at compilation time (the only dynamic components that are direct components are in this situation):



The record type GRAPH: the dynamic component X is a direct component.

The offset of an indirect component is always expressed in storage units.

The space reserved for the offset of an indirect component must be large enough to store the size of any value of the record type (the maximum potential offset). The compiler evaluates an upper bound MS of this size and treats an offset as a component having an anonymous integer type whose range is 0 .. MS.

If C is the name of an indirect component, then the offset of this component can be denoted in a component clause by the implementation generated name C'OFFSET.

4.8.3 Implicit Components

In some circumstances, access to an object of a record type or to its components involves computing information which only depends on the discriminant values. To avoid recomputation (which would degrade performance) the compiler stores this information in the record objects, updates it when the values of the discriminants are modified and uses it when the objects or its components are accessed. This information is stored in special components called implicit components.

An implicit component may contain information which is used when the record object or several of its components are accessed. In this case the component will be included in any record object (the implicit component is considered to be declared before any variant part in the record type declaration). There can be two components of this kind; one is called `RECORD_SIZE` and the other `VARIANT_INDEX`.

On the other hand an implicit component may be used to access a given record component. In that case the implicit component exists whenever the record component exists (the implicit component is considered to be declared at the same place as the record component). Components of this kind are called `ARRAY_DESCRIPTOR`s or `RECORD_DESCRIPTOR`s.

`'RECORD_SIZE`

This implicit component is created by the compiler when the record type has a variant part and its discriminants are defaulted. It contains the size of the storage space necessary to store the current value of the record object (note that the storage effectively allocated for the record object may be more than this).

The value of a `RECORD_SIZE` component may denote a number of bits or a number of storage units. In general it denotes a number of storage units, but if any component clause specifies that a component of the record type has an offset or a size which cannot be expressed using storage units, then the value designates a number of bits.

The implicit component `RECORD_SIZE` must be large enough to store the maximum size of any value of the record type. The compiler evaluates an upper bound `MS` of this size and then considers the implicit component as having an anonymous integer type whose range is `0 .. MS`.

If `R` is the name of the record type, this implicit component can be denoted in a component clause by the implementation generated name `R'RECORD_SIZE`. This allows user control over the position of the implicit component in the record.

`'VARIANT_INDEX`

This implicit component is created by the compiler when the record type has a variant part. It indicates the set of components that are present in a record value. It is used when a discriminant check is to be done.

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Component lists in variant parts that themselves do not contain a variant part are numbered. These numbers are the possible values of the implicit component VARIANT_INDEX.

```
type VEHICLE is (AIRCRAFT, ROCKET, BOAT, CAR);
```

```
type DESCRIPTION (KIND : VEHICLE := CAR) is
  record
```

```
    SPEED : INTEGER;
```

```
    case KIND is
```

```
      when AIRCRAFT | CAR =>
```

```
        WHEELS : INTEGER;
```

```
        case KIND is
```

```
          when AIRCRAFT =>      — 1
```

```
            WINGSPAN : INTEGER;
```

```
          when others => — 2
```

```
            null;
```

```
        end case;
```

```
      when BOAT => — 3
```

```
        STEAM : BOOLEAN;
```

```
      when ROCKET => — 4
```

```
        STAGES : INTEGER;
```

```
    end case;
```

```
  end record;
```

The value of the variant index indicates the set of components that are present in a record value:

Variant Index	Set
1	{KIND, SPEED, WHEELS, WINGSPAN}
2	{KIND, SPEED, WHEELS}
3	{KIND, SPEED, STEAM}
4	{KIND, SPEED, STAGES}

A comparison between the variant index of a record value and the bounds of an interval is enough to check that a given component is present in the value:

Component	Interval
KIND	—
SPEED	—
WHEELS	1 .. 2
WINGSPAN	1 .. 1
STEAM	3 .. 3
STAGES	4 .. 4

The implicit component `VARIANT INDEX` must be large enough to store the number `V` of component lists that don't contain variant parts. The compiler treats this implicit component as having an anonymous integer type whose range is `1 .. V`.

If `R` is the name of the record type, this implicit component can be denoted in a component clause by the implementation generated name `R'VARIANT INDEX`. This allows user control over the position of the implicit component in the record.

'ARRAY_DESCRIPTOR

An implicit component of this kind is associated by the compiler with each record component whose subtype is an anonymous array subtype that depends on a discriminant of the record. It contains information about the component subtype.

The structure of an implicit component of kind `ARRAY_DESCRIPTOR` is not described in this documentation. Nevertheless, if a programmer is interested in specifying the location of a component of this kind using a component clause, size of the component may be obtained using the `ASSEMBLY` parameter in the `COMPILE` command.

The compiler treats an implicit component of the kind `ARRAY_DESCRIPTOR` as having an anonymous array type. If `C` is the name of the record component whose subtype is described by the array descriptor, then this implicit component can be denoted in a component clause by the implementation generated name `C'ARRAY_DESCRIPTOR`. This allows user control over the position of the implicit component in the record.

'RECORD_DESCRIPTOR

An implicit component of this kind is associated by the compiler with each record component whose subtype is an anonymous record subtype that depends on a discriminant of the record. It contains information about the component subtype.

The structure of an implicit component of kind `RECORD_DESCRIPTOR` is not described in this documentation. Nevertheless, if a programmer is interested in specifying the location of a component of this kind using a component clause, the size of the component may be obtained using the `ASSEMBLY` parameter in the `COMPILE` command.

The compiler treats an implicit component of the kind `RECORD_DESCRIPTOR` as having an anonymous array type. If `C` is the name of the record component whose subtype is described by the record descriptor, then this implicit component can be denoted in a component clause by the implementation generated name `C'RECORD_DESCRIPTOR`. This allows user control over the position of the implicit component in the record.

Suppression of Implicit Components

The Alsys implementation provides the capability of suppressing the implicit

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components RECORD SIZE and/or VARIANT INDEX from a record type. This can be done using an implementation defined pragma called IMPROVE. The syntax of this pragma is as follows:

```
pragma IMPROVE ( TIME | SPACE , [ON =>] simple_name );
```

The first argument specifies whether TIME or SPACE is the primary criterion for the choice of the representation of the record type that is denoted by the second argument.

If TIME is specified, the compiler inserts implicit components as described above. If on the other hand SPACE is specified, the compiler only inserts a VARIANT INDEX or a RECORD SIZE component if this component appears in a record representation clause that applies to the record type. A record representation clause can thus be used to keep one implicit component while suppressing the other.

A pragma IMPROVE that applies to a given record type can occur anywhere that a representation clause is allowed for this type.

4.8.4 Size of Record Types and Objects

Size of a record subtype

Unless a component clause specifies that a component of a record type has an offset or a size which cannot be expressed using storage units, the size of a record subtype is rounded up to a whole number of storage units.

The size of a constrained record subtype is obtained by adding the sizes of its components and the sizes of its gaps (if any). This size is not computed at compile time

- o when the record subtype has non-static constraints,
- o when a component is an array or a record and its size is not computed at compile time.

The size of an unconstrained record subtype is obtained by adding the sizes of the components and the sizes of the gaps (if any) of its largest variant. If the size of a component or of a gap cannot be evaluated exactly at compile time an upper bound of this size is used by the compiler to compute the subtype size.

A size specification applied to a record type or first named subtype has no effect. The only size that can be specified using such a length clause is its usual size. Nevertheless, such a length clause can be useful to verify that the layout of a record is as expected by the application.

Size of an object of a record subtype

An object of a constrained record subtype has the same size as its subtype.

An object of an unconstrained record subtype has the same size as its subtype if this size is less than or equal to 8k bytes. If the size of the subtype is greater than this, the object has the size necessary to store its current value; storage space is allocated and released as the discriminants of the record change.

5 Conventions for Implementation-Generated Names

The following forms of implementation-generated names [13.4(8)] are used to denote implementation-dependent record components, as described in Section 4.8 in the sections on indirect and implicit components:

```
C'OFFSET
R'RECORD SIZE
R'VARIANT INDEX
R'ARRAY DESCRIPTORS
R'RECORD DESCRIPTORS
```

where C is the name of a record component and R the name of a record type.

The following predefined packages are reserved to Alsys and cannot be recompiled:

```
ALSYS_BASIC_IO
ALSYS_ADA_RUNTIME
ALSYS_BASIC_DIRECT_IO
ALSYS_BASIC_SEQUENTIAL_IO
```

6 Address Clauses

6.1 Address Clauses for Objects

An address clause can be used to specify an address for an object as described in RM 13.5. When such a clause applies to an object the compiler does not cause storage to be allocated for the object. The program accesses the object using the address specified in the clause. It is the responsibility of the user therefore to make sure that a valid allocation of storage has been done at the specified address.

An address clause is not allowed for task objects, for unconstrained records whose size is greater than 8k bytes or for a constant.

There are a number of ways to compose a legal address expression for use in an address clause. The most direct ways are:

- o For the case where the memory is defined in Ada as another object, use the 'ADDRESS attribute to obtain the argument for the address clause for the second object.
- o For the case where the desired location is memory defined in assembly or another non-Ada language (is relocatable), an interfaced routine may be

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used to obtain the appropriate address from referencing information known to the other language.

- o For the case where an absolute address is known to the programmer, use the function `SYSTEM.VALUE`. This function has one (1) parameter of type string. The string is of the form "ssss:0000". The length is 9, and the ':' is required to separate the segment and offset portion of the address. `SYSTEM.VALUE` returns a value of type `SYSTEM.ADDRESS`. The string representing the desired address can be passed as the actual parameter to `SYSTEM.VALUE` in the simple expression part of the address clause. See Section 3 for the specification of package `SYSTEM`.

In all cases other than the use of an address attribute, the programmer must ensure that the segment part of the argument is a selector if the program is to run in protected mode. Refer to the Application Developers' Guide, Section 5.1.5 for more information on protected mode machine oriented programming.

6.2 Address Clauses for Program Units

Address clauses for program units are not implemented in the current version of the compiler.

6.3 Address Clauses for Interrupt Entries

Address clauses for interrupt entries are supported. (See Chapter 7 of the Application Developer's Guide for details.)

7 Unchecked Conversions

Unchecked type conversions are described in [RM 13.10.2]. The following restrictions apply to their use:

- o Unconstrained arrays are not allowed as target types. Unconstrained record types without defaulted discriminants are not allowed as target types. Access types to unconstrained arrays are not allowed as target or source types. Notes also that `UNCHECKED CONVERSION` cannot be used for an access to an unconstrained string.
- o If the target type has a smaller size than the

source type then the target is made of the least significant bits of the source.

If the source and the target types are each of scalar or access type or if they are both of composite type, the effect of the function is to return the operand.

In other cases the effect of unchecked conversion can be considered as a copy:

- o If an unchecked conversion of a scalar or access source type to a composite target type is performed, the result is a copy of the source operand. The result has the size of the source.
- o If an unchecked conversion of a composite source type to a scalar or access target type is performed, the result is a copy of the source operand. The result has the size of the target.

8 Input-Output Packages

The RM defines the predefined input-output package TEXT IO, and describes how to use the facilities available within this packages. The RM also defines the package IO EXCEPTIONS, which specifies the exceptions that can be raised by the predefined input-output package.

The only input-output supported is TEXT IO to standard-input and standard-output. The predefined packages SEQUENTIAL IO and DIRECT IO are not supported. No file input-output is support at all. Any attempted file operations (such as OPEN, CLOSE, CREATE, or DELETE) will cause a USE_ERROR to be generated.

In addition the RM outlines the package LOW_LEVEL IO, which is concerned with low-level machine-dependent input-output, such as would possibly be used to write device drivers or access device registers. LOW_LEVEL IO has not been implemented. However the user provides low level drivers to support a specific hardware.

Sections 8.1 through 8.4 apply to I/O in cross mode. Sections 8.5 through 8.14 apply to I/O in native mode.

8.1 Accessing Devices

The user must provide a description of the input-output devices and provide drivers for them. These drivers will permit the various devices to be accessed through the standard I/O package TEXT IO.

All necessary information to describe the devices and build the drivers is available in the Cross Development Guide.

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8.2 The Form Function

The function FORM always returns a null string.

8.3 Text Input-Output

All text column numbers, line numbers, and page numbers are values of the subtype POSITIVE_COUNT.

If STANDARD_INPUT and STANDARD_OUTPUT are the console input and output devices (as indicated by the function ADA@IO_IOCTL), prompts written to STANDARD_OUTPUT with the procedure PUT will appear before (or when) a GET (or GET_LINE) occurs.

The functions END_OF_PAGE and END_OF_FILE always return FALSE as input-output is only performed for a device.

END_ERROR is raised for STANDARD_INPUT when CTRL-Z (ASCII.SUB) is entered at the console input device.

8.4 Tasking Issues

Care should be taken when performing multiple input-output operations during tasking because the order of calls to the I/O primitives is unpredictable. For example, two strings output by TEXT_IO.PUT_LINE in two different tasks may appear with interleaved characters. Synchronization of I/O in cases such as this is the user's responsibility.

The TEXT_IO files STANDARD_INPUT and STANDARD_OUTPUT are shared by all tasks of an Ada program.

If TEXT_IO.STANDARD_INPUT corresponds to the console input device, it will not block a program on input. All tasks not waiting for input will continue running.

The following Sections 8.5 - 8.14 apply to native mode only.

8.5 Correspondence between External Files and DOS Files

Ada input-output is defined in terms of external files. Data is read from and written to external files. Each external file is implemented as a standard DOS file, including the use of STANDARD_INPUT and STANDARD_OUTPUT.

The name of an external file can be either

- o the null string
- o a DOS filename
- o a DOS special file or device name (for example, CON and PRN)

If the name is a null string, the associated external file is a temporary file and will cease to exist when the program is terminated. The file will be placed in the current directory and its name will be chosen by DOS.

If the name is a DOS filename, the filename will be interpreted according to standard DOS conventions (that is, relative to the current directory). The exception NAME_ERROR is raised if the name part of the filename has more than 8 characters or if the extension part has more than 3 characters.

If an existing DOS file is specified to the CREATE procedure, the contents of the file will be deleted before writing to the file.

If a non-existing directory is specified in a filepath name to CREATE, the directory will not be created, and the exception NAME_ERROR is raised.

8.6 Error Handling

DOS errors are translated into Ada exceptions, as defined in the RM by package IO_EXCEPTIONS. In particular, DEVICE_ERROR is raised in cases of drive not ready, unknown media, disk full or hardware errors on the disk (such as read or write fault).

8.7 The FORM Parameter

The form parameter is a string, formed from a list of attributes, with attributes separated by commas. The string is not case sensitive. The attributes specify:

- o Buffering

BUFFER_SIZE => size_in_bytes

- o Appending

APPEND => YES | NO

- o Truncation of the name by DOS

TRUNCATE => YES | NO

- o DIRECT_IO on UNCONSTRAINED objects

RECORD_SIZE => size_in_bytes

where:

BUFFER_SIZE: Controls the size of the internal buffer. This option is not supported for DIRECT_IO. The default value is 1024. This option has no effect when used by TEXT_IO with an external file that is a character device, in which case the size of the buffer will be 0.

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APPEND: If YES output is appended to the end of the existing file. If NO output overwrites the existing file. This option is not supported for `DIRECT_IO`. The default is NO.

TRUNCATE: If YES the file name will be automatically truncated if it is bigger than 8 characters. The default value is NO, meaning that the exception `NAME_ERROR` will be raised if the name is too long.

RECORD SIZE: This option is supported only for `DIRECT_IO`. This attribute controls the logical record length of the external file.

- When `DIRECT_IO` is instantiated with an unconstrained type the user is required to specify the `RECORD SIZE` attribute (otherwise `USE_ERROR` will be raised). The value given must be larger or equal to the largest record which is going to be written. If a larger record is processed the exception `USE_ERROR` will be raised.
- When `DIRECT_IO` is instantiated with a constrained type the user is not required to specify the `RECORD SIZE` but if the `RECORD SIZE` is specified the only possible value would be the element size in bytes. Any other values will raise `USE_ERROR`.

The exception `USE_ERROR` is raised if the form `STRING` is not correct or if a non supported attribute for a given package is used.

Example:

```
FORM => "TRUNCATE => YES, APPEND => YES, BUFFER_SIZE => 20480"
```

8.8 Sequential Files

For sequential access the file is viewed as a sequence of values that are transferred in the order of their appearance (as produced by the program or run-time environment). This is sometimes called a stream file in other operating systems. Each object in a sequential file has the same binary representation as the Ada object in the executable program.

8.9 Direct Files

For direct access the file is viewed as a set of elements occupying consecutive

positions in a linear order. The position of an element in a direct file is specified by its index, which is an integer of subtype `POSITIVE_COUNT`.

`DIRECT IO` only allows input-output for constrained types. If `DIRECT IO` is instantiated for an unconstrained type, all calls to `CREATE` or `OPEN` will raise `USE_ERROR`. Each object in a direct file will have the same binary representation as the Ada object in the executable program. All elements within the file will have the same length.

8.10 Text Files

Text files are used for the input and output of information in ASCII character form. Each text file is a sequence of characters grouped into lines, and lines are grouped into a sequence of pages.

All text file column numbers, line numbers, and page numbers are values of the subtype `POSITIVE_COUNT`.

Note that due to the definitions of line terminator, page terminator, and file terminator in the RM, and the method used to mark the end of file under DOS, some ASCII files do not represent well-formed `TEXT_IO` files.

A text file is buffered by the Runtime Executive unless

- o it names a device (for example, `CON` or `PRN`).
- o it is `STANDARD_INPUT` or `STANDARD_OUTPUT` band has not been redirected.

If not redirected, prompts written to `STANDARD_OUTPUT` with the procedure `PUT` will appear before (or when) a `GET` (or `GET_LINE`) occurs.

The functions `END OF PAGE` and `END OF FILE` always return `FALSE` when the file is a device, which includes the use of the file `CON`, and `STANDARD_INPUT` when it is not redirected. Programs which would like to check for end of file when the file may be a device should handle the exception `END_ERROR` instead, as in the following example:

Example

```
begin
  loop
    -- Display the prompt:
    TEXT_IO.PUT ("—> ");
    -- Read the next line:
    TEXT_IO.GET_LINE (COMMAND, LAST);
    -- Now do something with COMMAND (1 .. LAST)
  end loop;
exception
  when TEXT_IO.END_ERROR =>
    null;
end;
```

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END_ERROR is raised for STANDARD_INPUT when CTRL-Z (ASCII.SUB) is entered at the keyboard.

8.11 Access Protection of External Files

All DOS access protections exist when using files under DOS. If a file is open for read only access by one process it can not be opened by another process for read/write access.

8.12 The Need to Close a File Explicitly

The Runtime Executive will flush all buffers and close all open files when the program is terminated, either normally or through some exception.

However, the RM does not define what happens when a program terminates without closing all the opened files. Thus a program which depends on this feature of the Runtime Executive might have problems when ported to another system.

8.13 Limitation on the Procedure RESET

An internal file opened for input cannot be RESET for output. However, an internal file opened for output can be RESET for input, and can subsequently be RESET back to output.

8.14 Sharing of External Files and Tasking Issues

Several internal files can be associated with the same external file only if all the internal files are opened with mode IN_MODE. However, if a file is opened with mode OUT_MODE and then changed to IN_MODE with the RESET procedure, it cannot be shared.

Care should be taken when performing multiple input-output operations on an external file during tasking because the order of calls to the I/O primitives is unpredictable. For example, two strings output by TEXT_IO.PUT_LINE in two different tasks may appear in the output file with interleaved characters. Synchronization of I/O in cases such as this is the user's responsibility.

The TEXT_IO files STANDARD_INPUT and STANDARD_OUTPUT are shared by all tasks of an Ada program.

If TEXT_IO.STANDARD_INPUT is not redirected, it will not block a program on input. All tasks not waiting for input will continue running.

9 Characteristics of Numeric Types

9.1 Integer Types

The ranges of values for integer types declared in package STANDARD are as follows:

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SHORT_INTEGER	-128 .. 127	— $2^{**7} - 1$
INTEGER	-32768 .. 32767	-- $2^{**15} - 1$
LONG_INTEGER	-2147483648 .. 2147483647	— $2^{**31} - 1$

For the packages `DIRECT_IO*` and `TEXT_IO`, the range of values for types `COUNT` and `POSITIVE_COUNT` are as follows:

COUNT	0 .. 2147483647	— $2^{**31} - 1$
POSITIVE_COUNT	1 .. 2147483647	— $2^{**31} - 1$

For the package `TEXT_IO`, the range of values for the type `FIELD` is as follows:

FIELD	0 .. 255	— $2^{**8} - 1$
-------	----------	-----------------

* Note: `DIRECT_IO` only applies to native mode.

9.2 Floating Point Type Attributes

	SHORT FLOAT and FLOAT	LONG_FLOAT
DIGITS	6	15
MANTISSA	21	51
EMAX	84	204
EPSILON	9.53674E-07	8.88178E-16
LARGE	1.93428E+25	2.57110E+61
SAFE_EMAX	125	1021
SAFE_SMALL	1.17549E-38	2.22507E-308
SAFE_LARGE	4.25353E+37	2.24712E+307
FIRST	-3.40282E+38	-1.79769E+308
LAST	3.40282E+38	1.79769E+308
MACHINE_RADIX	2	2
MACHINE_MANTISSA	24	53

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MACHINE_EMAX	128	1024
MACHINE_EMIN	-125	-1021
MACHINE_ROUNDS	true	true
MACHINE_OVERFLOWS	false	false
SIZE	32	64

9.3 Attributes of Type DURATION

DURATION'DELTA	2.0 ** (-14)
DURATION'SMALL	2.0 ** (-14)
DURATION'FIRST	-131_072.0
DURATION'LAST	131_072.0
DURATION'LARGE	same as DURATION'LAST

10 Other Implementation-Dependent Characteristics

10.1 Use of the Floating-Point Coprocessor

Floating point coprocessor instructions are used in programs that perform arithmetic on floating point values in some fixed point operations and when the FLOAT_IO or FIXED_IO packages of TEXT_IO are used. The mantissa of a fixed point value may be obtained through a conversion to an appropriate integer type. This conversion does not use floating point operations. Object code running on an 80286 or 80386 using floating point instructions does not require the coprocessor, since software floating point emulation is provided (see Binder option FLOAT in User's Guide, Section 5.2). Object code running on an 8086, 8088 or 80186 does require an 8087 coprocessor, since 8087 software emulation is not supported. See Appendix D of the Application Developer's Guide for more details.

The Runtime Executive will detect the absence of the floating point coprocessor if it is required by a program and will raise CONSTRAINT_ERROR.

10.2 Characteristics of the Heap

All objects created by allocators go into the heap. Also, portions of the Runtime Executive representation of task objects, including the task stacks, are allocated in the heap.

UNCHECKED DEALLOCATION is implemented for all Ada access objects except access objects to tasks. Use of UNCHECKED DEALLOCATION on a task object will lead to unpredictable results.

All objects whose visibility is linked to a subprogram, task body, or block have their storage reclaimed at exit, whether the exit is normal or due to an exception. Effectively pragma CONTROLLED is automatically applied to all access types. Moreover, all compiler temporaries on the heap (generated by such operations as function calls returning unconstrained arrays, or many concatenations) allocated in a scope are deallocated upon leaving the scope.

Note that the programmer may force heap reclamation of temporaries associated with any statements by enclosing the statement in a begin .. end block. This is especially useful when complex concatenations or other heap-intensive operations are performed in loops, and can reduce or eliminate STORAGE_ERRORS that might otherwise occur.

The maximum size of the heap is limited only by available memory. This includes the amount of physical memory (RAM) and the amount of virtual memory (hard disk swap space).

10.3 Characteristics of Tasks

The default task stack size is 1K bytes (32K bytes for the environment task), but by using the Binder option STACK.TASK the size for all task stacks in a program may be set to a size from 1K bytes to 64K bytes.

Normal priority rules are followed for preemption, where PRIORITY values are in the range 1 .. 10. A task with undefined priority (no pragma PRIORITY) is considered to be lower than priority 1.

The maximum number of active tasks is restricted only by memory usage.

The acceptor of a rendezvous executes the accept body code in its own stack. Rendezvous with an empty accept body (for synchronization) does not cause a context switch.

The main program waits for completion of all tasks dependent upon library packages before terminating.

Abnormal completion of an aborted task takes place immediately, except when the abnormal task is the caller of an entry that is engaged in a rendezvous, or if it is in the process of activating some tasks. Any such task becomes abnormally completed as soon as the state in question is exited.

The message

GLOBAL BLOCKING SITUATION DETECTED

is printed to STANDARD_OUTPUT when the Runtime Executive detects that no further progress is possible for any task in the program. The execution of the program is then abandoned.

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10.4 Definition of a Main Subprogram

A library unit can be used as a main subprogram if and only if it is a procedure that is not generic and that has no formal parameters.

10.5 Ordering of Compilation Units

The compiler imposes no additional ordering constraints on compilations beyond those required by the language.

11 Limitations

11.1 Compiler Limitations

- o The maximum identifier length is 255 characters.
- o The maximum line length is 255 characters.
- o The maximum number of unique identifiers per compilation unit is 2500.
- o The maximum number of compilation units in a library is 1000.
- o The maximum number of Ada libraries in a family is 15.
- o If an unconstrained record type can exceed 4096 bytes, the type is not permitted (unless constrained) as the element type in the definition of an array or record type.

11.2 Hardware Related Limitations

- o The maximum amount of data in the heap is limited only by available memory.
- o The maximum size of the generated code for a single compilation unit is 65535 bytes.
- o The maximum size of a single array or record object is 65522 bytes. An object bigger than 4096 bytes will be indirectly allocated. Refer to ALLOCATION parameter in the COMPILE command. (Section 4.2 of the User's Guide.)
- o The maximum size of a single stack frame is 32766 bytes, including the data for inner package subunits unnested to the parent frame.
- o The maximum amount of data in the global data area is 65535 bytes, including compiler generated data that goes into the GDA (about 8 bytes per compilation unit plus 4 bytes per externally visible subprogram).